

Simplicity of Nekrashevych algebras of contracting self-similar groups

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Cimpa School — From Dynamics to Algebra and
Representation Theory and Back

Outline

Self-similar groups

Nekrashevych algebras

Techniques

Rooted trees

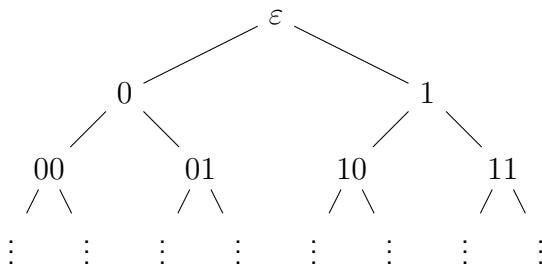
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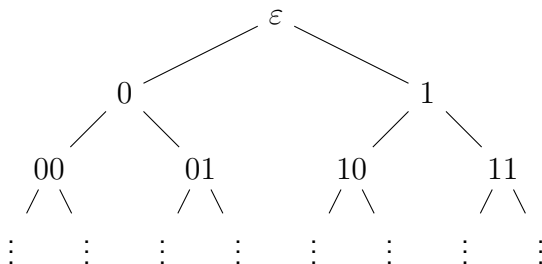
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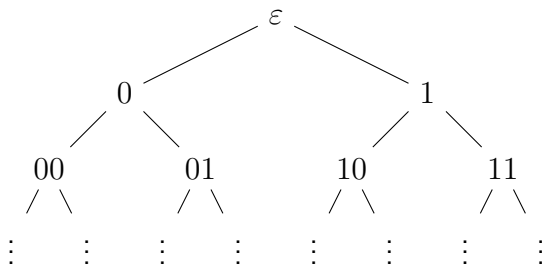
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- So $\text{Aut}(T_X)$ acts on X^ω .

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- Note that $g|_{x_1 \dots x_n} = ((g|_{x_1})|_{x_2} \dots)|_{x_n}$.

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- Bartholdi and Nekrashevych used self-similar groups called iterated monodromy groups to solve Hubbard's twisted rabbit problem in complex dynamics.
- Any self-similar subgroup of $\text{Aut}(T_X)$ acts on X^ω .

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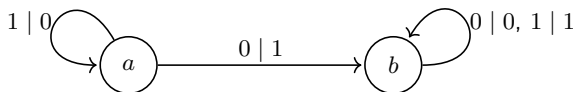
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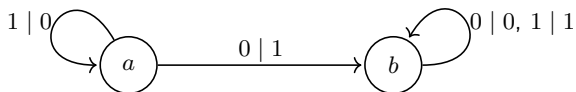
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- Automaton groups always have decidable word problem.

Adding machine



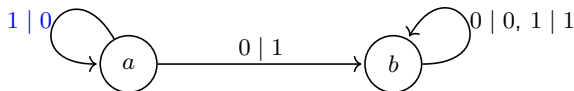
- $a|_0 = b, a|_1 = a, a(0) = 1, a(1) = 0.$

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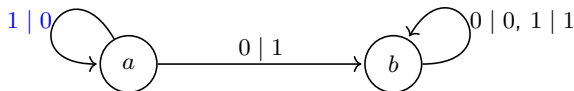
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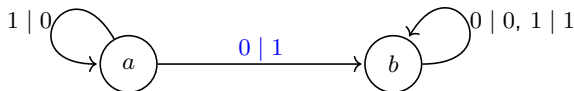
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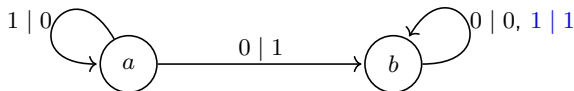
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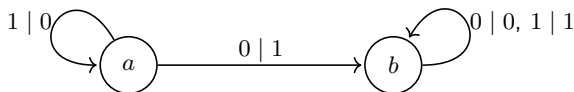
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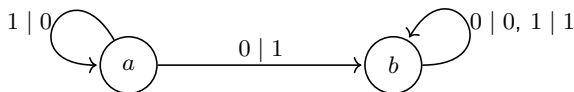
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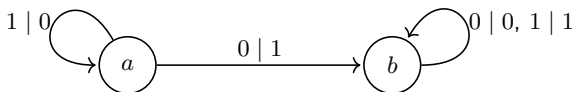
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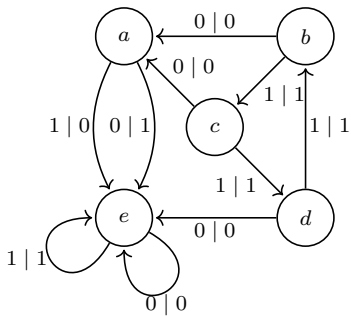
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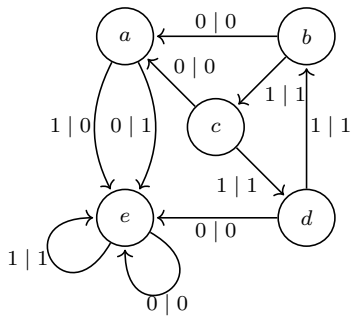


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- $\langle a, b \rangle \cong \mathbb{Z}$.

The Grigorchuk group

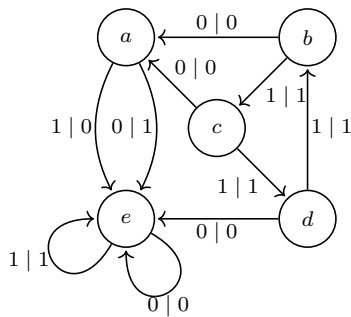


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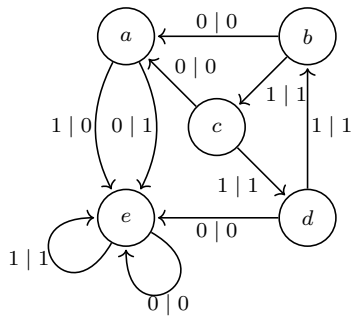
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- It was the first amenable but not elementarily amenable group found.

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- Sections of g along sufficiently long words will then have small word length, allowing induction.
- Nekrashevych formalized the notion in a way that makes sense for any self-similar group.

The nucleus

- A self-similar group $G \leq \text{Aut}(T_X)$ is **contracting** if there is a finite automaton $\mathcal{N} \subseteq G$ such that, for all $g \in G$, there is $k \geq 0$ with $g|_{X^k} \subseteq \mathcal{N}$.

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- Contracting groups include the Grigorchuk group, Gupta-Sidki groups, GGS-groups, Šunić groups associated to polynomials, automaton spinal groups, the basilica group and the Hanoi towers group on 3-pegs.

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- The C^* -algebra is a completion of the complex Nekrashevych algebra.
- The algebraic version generalizes Leavitt algebras.
- They have also been studied by Clark, Exel, Pardo, Sims, Starling, the authors and others.

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- It is a finitely presented simple K -algebra acting faithfully on KX^ω .

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- The **Nekrashevych algebra** $N_K(G, X)$ of G is the K -algebra generated by G and the x, x^* with $x \in X$ subject to the Leavitt algebra relations and the above relations.

The simple quotient

Theorem (BS, Szakács)

The K -algebra generated by $L_{K,X}$ and G acting on KX^ω is the unique simple quotient of the Nekrashevych algebra of G .

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- So simplicity of $N_K(G, X)$ is equivalent to faithfulness of the natural representation on infinite words.

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- But these groupoids are rarely Hausdorff, so they are good test cases for understanding simplicity phenomena.

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- Nekrashevych (unpublished) showed the Nekrashevych algebra of the Grigorchuk-Erschler group is simple over no field.

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- There is a theoretical description of simplicity that underlies the algorithmic result.
 - This theoretical description is applicable to many infinite families of contracting groups.
 - For any finite set \mathcal{P} of primes, we construct a contracting group G with $N_K(G, X)$ simple over precisely fields K of characteristic not in \mathcal{P} .

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- Šunić gave necessary and sufficient conditions for $G_{p,f}$ to be a p -group of intermediate growth.

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Theorem (BS, Szakács)

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 - The Grigorchuk-Erschler group G_{2,x^2+1} has a non-simple algebra because the companion matrix fixes $(1, 1)$.

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- We also gave characterizations of simplicity for self-replicating spinal automaton groups and multi-edge spinal groups.

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- The ideal $(1 - \sum_{x \in X} xx^*)$ is a special case of Exel's tight ideal of an inverse monoid algebra.

Theorem (BS, Szakács)

*The algebra of a congruence-free inverse monoid has a unique maximal ideal containing Exel's tight ideal, called the **singular ideal**.*

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 - If G is contracting, we may assume that a is supported on the nucleus \mathcal{N} , a finite set.

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- The obstruction to solving the general case is bounding the max word length in the support of a singular element of KG not belonging to $(1 - \sum_{x \in X} xx^*)$.

The end

Thank you for your attention!